

A Game Theoretic Model for User Preference-Aware Resource Pricing in Wireless Mobile Networks

Undergraduate Research Report

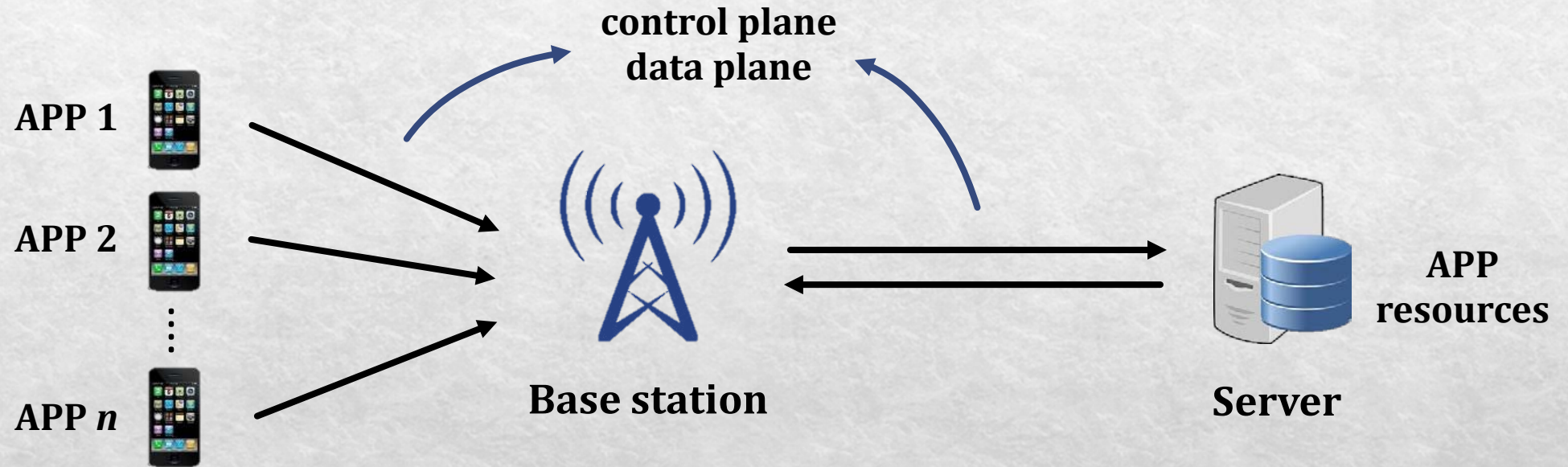
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Outline

- ◇ Game Model
- ◇ Solving Stackelberg Game
- ◇ Results

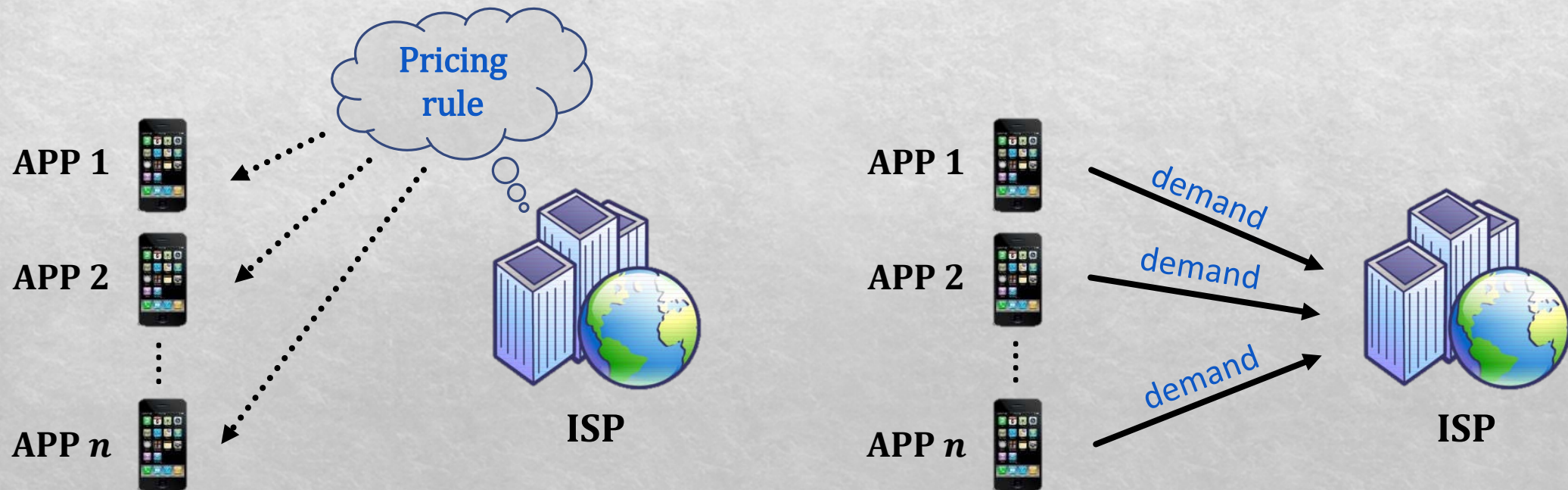
Game Model



- ◇ Players: APP 1, APP 2, ... , APP n and ISP
- ◇ Resources: R_C (Control), R_D (Data), R_A (APP)

Non-cooperative Stackelberg Game

- ◆ First the ISP announces a pricing rule for the users.
- ◆ Then the users request for resource based on their demand.



Utility Functions

For user i :

$$U_i = \min \left\{ w_i \log \left(1 + \frac{x_{C,i}}{a_i} \right), w_i \log \left(1 + \frac{x_{D,i}}{b_i} \right), w_i \log \left(1 + \frac{x_{A,i}}{c_i} \right) \right\} \\ - (p_C x_{C,i} + p_D x_{D,i} + p_A x_{A,i})$$

For the ISP :

$$U_{ISP} = \sum_i (p_C x_{C,i} + p_D x_{D,i} + p_A x_{A,i})$$

p_C, p_D, p_A : per data price $x_{C,i}, x_{D,i}, x_{A,i}$: demand of user i

w_i : user types a_i, b_i, c_i : user preference, where $a_i + b_i + c_i = 1$

Two-stage Stackelberg Game

$$\begin{cases} U_i = \min \left\{ w_i \log \left(1 + \frac{x_{C,i}}{a_i} \right), w_i \log \left(1 + \frac{x_{D,i}}{b_i} \right), w_i \log \left(1 + \frac{x_{A,i}}{c_i} \right) \right\} \\ \quad - (p_C x_{C,i} + p_D x_{D,i} + p_A x_{A,i}) \\ U_{ISP} = \sum_i (p_C x_{C,i} + p_D x_{D,i} + p_A x_{A,i}) \end{cases}$$

Stage 1 – Given p_C, p_D, p_A , maximize U_i

Stage 2 – Given $(x_{C,i}^*, x_{D,i}^*, x_{A,i}^*) = f(p_C, p_D, p_A)$, maximize U_{ISP}

Solving Stackelberg Game

Step 1 - Given p_C, p_D, p_A , maximize U_i

$$U_i = \min \left\{ w_i \log \left(1 + \frac{x_{C,i}}{a_i} \right), w_i \log \left(1 + \frac{x_{D,i}}{b_i} \right), w_i \log \left(1 + \frac{x_{A,i}}{c_i} \right) \right\} \\ - (p_C x_{C,i} + p_D x_{D,i} + p_A x_{A,i})$$

(1) From the minimal function part, we get

$$w_i \log \left(1 + \frac{x_{C,i}}{a_i} \right) = w_i \log \left(1 + \frac{x_{D,i}}{b_i} \right) = w_i \log \left(1 + \frac{x_{A,i}}{c_i} \right) \Rightarrow \frac{x_{C,i}}{a_i} = \frac{x_{D,i}}{b_i} = \frac{x_{A,i}}{c_i}$$

Solving Stackelberg Game

Step 1 - Given p_C, p_D, p_A , maximize U_i

$$U_i = w_i \log \left(1 + \frac{x_{C,i}}{a_i} \right) - (p_C x_{C,i} + p_D x_{D,i} + p_A x_{A,i})$$

(2) Now we maximize U_i

$$\frac{\partial U_{User\ i}}{\partial x_{C,i}} = \frac{\partial U_{User\ i}}{\partial x_{D,i}} = \frac{\partial U_{User\ i}}{\partial x_{A,i}} = 0 \Rightarrow (x_{C,i}, x_{D,i}, x_{A,i}) = (a_i k_i, b_i k_i, c_i k_i)$$

$$\text{where } k_i = \frac{w_i}{a_i p_C + b_i p_D + c_i p_A} - 1$$

Solving Stackelberg Game

Step 2 - Given $(x_{C,i}^*, x_{D,i}^*, x_{A,i}^*) = f(p_C, p_D, p_A)$, maximize U_{ISP}

$$U_{ISP} = \sum_i (p_C x_{C,i} + p_D x_{D,i} + p_A x_{A,i})$$

$$\Rightarrow \frac{\partial U_{ISP}}{\partial p_C} < 0 \forall p_C > 0, \quad \frac{\partial U_{ISP}}{\partial p_D} < 0 \forall p_D > 0, \quad \frac{\partial U_{ISP}}{\partial p_A} < 0 \forall p_A > 0$$

Capacity Limitation:

$$\sum_i (x_{C,i} + x_{D,i} + x_{A,i}) \leq C \Leftrightarrow \sum_i \left(\frac{w_i}{a_i p_C + b_i p_D + c_i p_A} \right) \leq C + N$$

Solving Stackelberg Game

Step 2 - Given $(x_{C,i}^*, x_{D,i}^*, x_{A,i}^*) = f(p_C, p_D, p_A)$, maximize U_{ISP}

Goal:

$$\text{maximize } U_{ISP} = \sum_i (p_C x_{C,i} + p_D x_{D,i} + p_A x_{A,i})$$

$$\text{with } \sum_i \left(\frac{w_i}{a_i p_C + b_i p_D + c_i p_A} \right) \leq C + N$$

Solving Stackelberg Game

Step 2 - Given $(x_{C,i}^*, x_{D,i}^*, x_{A,i}^*) = f(p_C, p_D, p_A)$, maximize U_{ISP}

Capacity Limitation Plane: an approximated border of feasible region

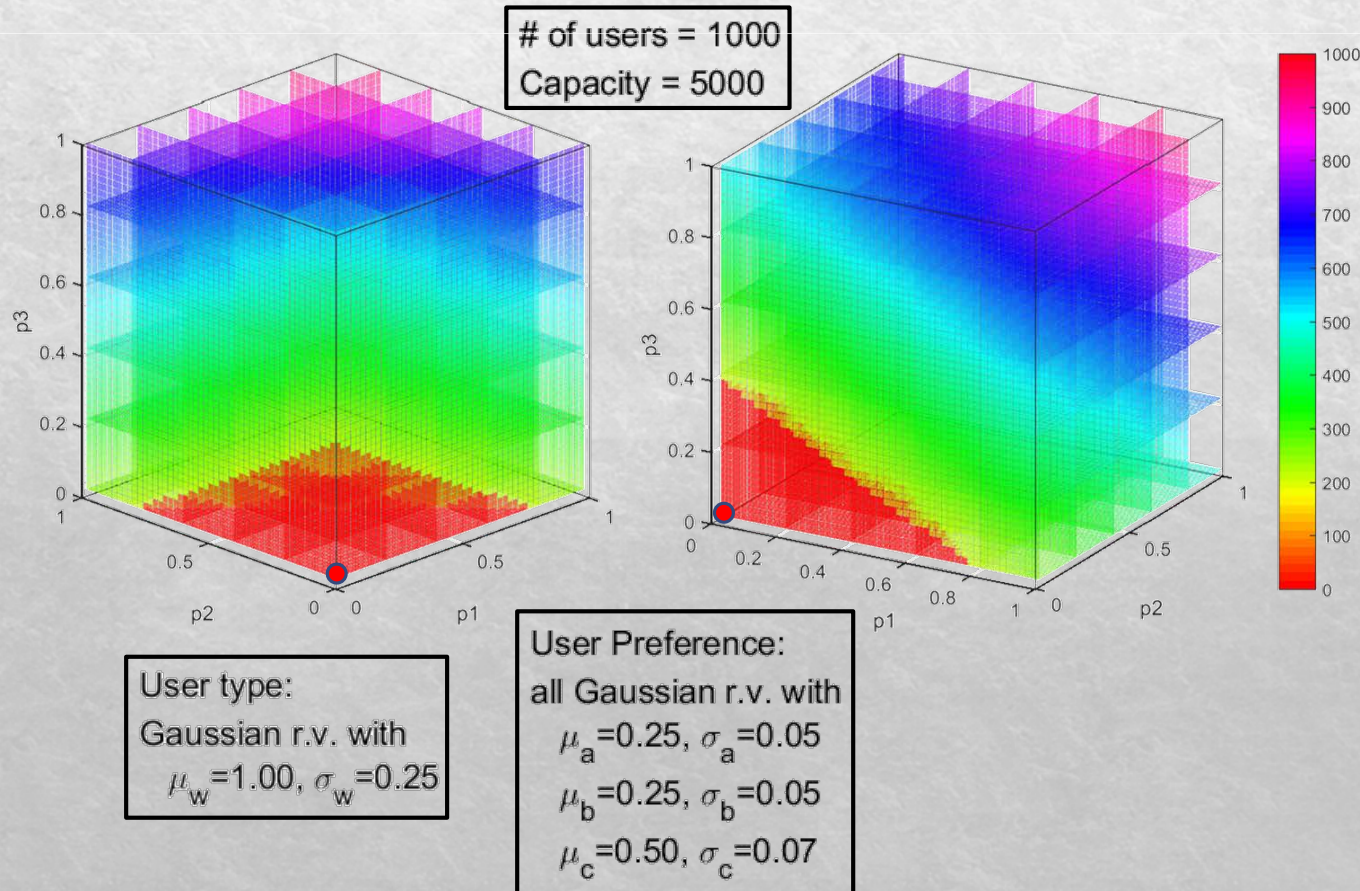
$$\sum_i \left(\frac{w_i}{a_i p_C + b_i p_D + c_i p_A} \right) = C + N$$

assuming small variance on a_i, b_i, c_i , we get

$$\Rightarrow \mu_a p_C + \mu_b p_D + \mu_c p_A = \frac{\sum w_i}{C + N}$$

Solving Stackelberg Game

Step 2 - Given $(x_{C,i}^*, x_{D,i}^*, x_{A,i}^*) = f(p_C, p_D, p_A)$, maximize U_{ISP}



◇ Generally, the utility of ISP increase as the per data price goes down.

◇ The maximum revenue of ISP can be achieved around the capacity limitation plane:
$$\sum_i (x_{C,i} + x_{D,i} + x_{A,i}) = C.$$

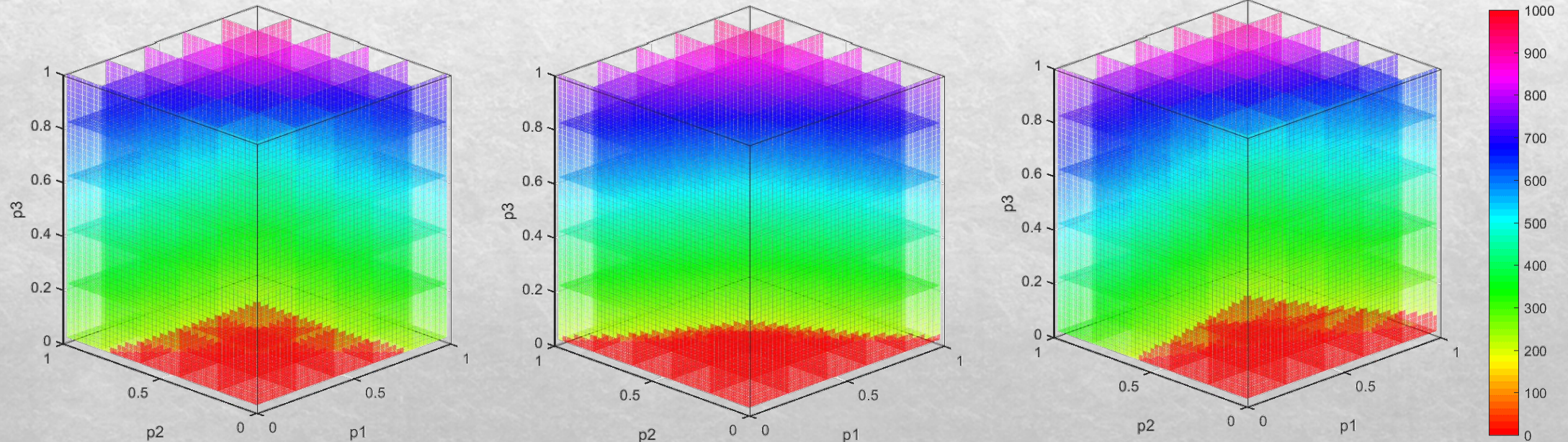
Solving Stackelberg Game

The limitation plane's direction changes with different combinations of user preference.

Effect of user preference

of users = 1000
Capacity = 5000

User type:
Gaussian r.v. with
 $\mu_w = 1.00, \sigma_w = 0.25$



$$(\mu_a, \mu_b, \mu_c) = (0.25, 0.25, 0.5)$$

$$(\mu_a, \mu_b, \mu_c) = (0.2, 0.2, 0.6)$$

$$(\mu_a, \mu_b, \mu_c) = (0.17, 0.33, 0.5)$$

Solving Stackelberg Game

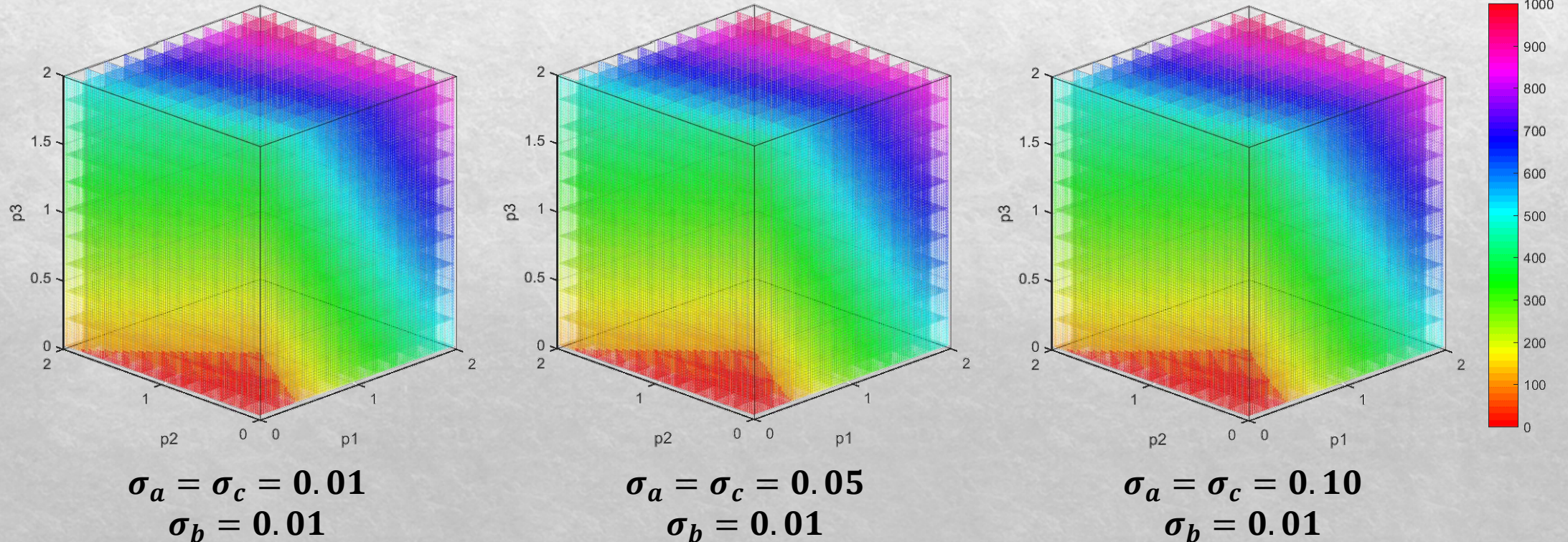
The border of feasible region becomes rough when the variance of user type increase.

Effect of user type variance

of users = 1000
Capacity = 5000

$$\sigma_a = \sigma_c = 0.01$$
$$\sigma_b = 0.01$$

$$(\mu_a, \mu_b, \mu_c) = (0.5, 0.1, 0.4)$$



Solving Stackelberg Game

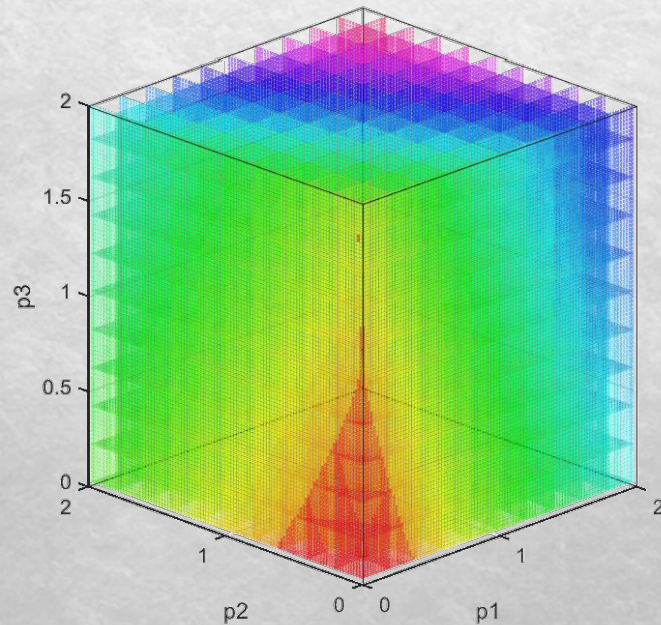
The border of feasible region becomes rough when the variance of user type increase.

Effect of user type variance

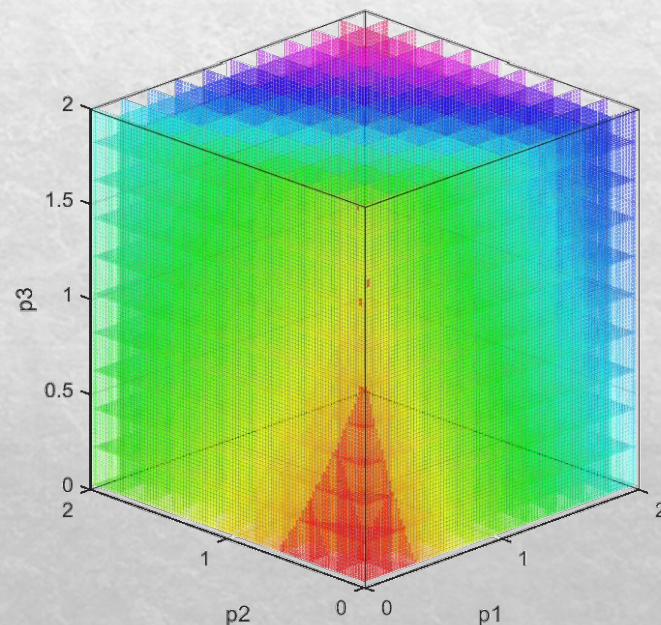
of users = 1000
Capacity = 5000

User type:
Gaussian r.v. with
 $\mu_w = 1.00$, $\sigma_w = 0.25$

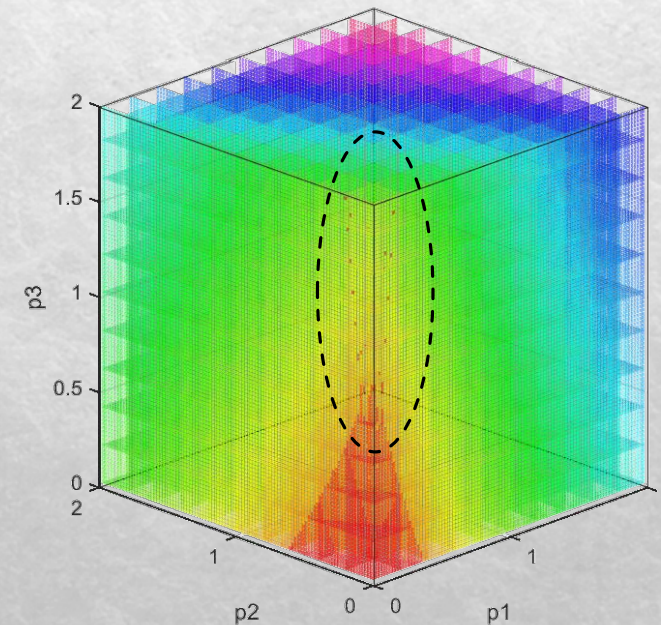
$(\mu_a, \mu_b, \mu_c) = (0.5, 0.3, 0.2)$
 $(\sigma_a, \sigma_b, \sigma_c) = (0.05, 0.05, 0.07)$



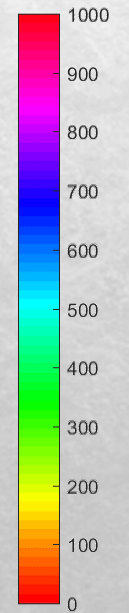
$\sigma_w = 0.05$



$\sigma_w = 0.15$



$\sigma_w = 0.25$



Results

- ◆ The optimal combination of (p_C, p_D, p_A) that results in the maximum revenue of the ISP is around **Capacity Limitation Plane**:

$$\mu_a p_C + \mu_b p_D + \mu_c p_A = \frac{\sum w_i}{C + N}$$

- ◆ Picking any combination of (p_C, p_D, p_A) on **Capacity Limitation Plane** results in an approximately optimal utility of the ISP.
- ◆ If we want to simplify the pricing rule, we can simply choose

$$p_C = p_D = p_A = \frac{\sum w_i}{3\mu(C + N)}$$

Thank you for your attention!

Appendix

Lemma

Given a random variable X with mean μ_x and variance σ_x ,

$$E[(x - \mu_x)^n] \ll \mu_x^n \quad \forall n > 2 \quad \Rightarrow \quad E\left[\frac{1}{X}\right] \sim \frac{1}{\mu_x} + \frac{\sigma_x^2}{2\mu_x^2}$$

Proof:

$$g(X) = g(\mu_x + X - \mu_x) = g(\mu_x) + (X - \mu_x)g'(\mu_x) + \frac{(X - \mu_x)^2}{2!}g''(\mu_x) + \dots$$

$$E[g(X)] = E[g(\mu_x)] + E[(X - \mu_x)g'(\mu_x)] + E\left[\frac{(X - \mu_x)^2}{2!}g''(\mu_x)\right] + \dots$$

Appendix

Proof: (continued)

$$\begin{aligned} E[g(X)] &= E[g(\mu_x)] + E[(X - \mu_x)g'(\mu_x)] + E\left[\frac{(X - \mu_x)^2}{2!}g''(\mu_x)\right] + \dots \\ &= g(\mu_x) + 0 + \frac{\sigma_x^2}{2}g''(\mu_x) + \frac{E[(x - \mu_x)^n]}{n!}g^{(n-1)}(\mu_x) \\ &= \mu_x + 0 + \frac{\sigma_x^2}{2}g''(\mu_x) + \frac{E[(x - \mu_x)^n]}{n!}g^{(n-1)}(\mu_x) \\ &= \mu_x + 0 + \frac{\sigma_x^2}{12\mu_x^3} + \frac{1}{n!} \frac{E[(x - \mu_x)^n]}{(n-1)!\mu_x^n} \\ &\sim \mu_x + 0 + \frac{\sigma_x^2}{12\mu_x^3} \quad \text{by } E[(x - \mu_x)^n] \ll \mu_x^n \end{aligned}$$

Appendix

In our model, consider a_i, b_i, c_i as random variables A, B, C, respectively.

Let $x_i = a_i p_C + b_i p_D + c_i p_A$, then

$$\mu_x = \mu_a p_C + \mu_b p_D + \mu_c p_A, \quad \sigma_x = \sigma_a p_C + \sigma_b p_D + \sigma_c p_A$$

If $\forall n > 2, E[(x_i - \mu_x)^n] \ll \mu_x^n$, then we have

$$E \left[\frac{1}{a_i p_C + b_i p_D + c_i p_A} \right] \sim \frac{1}{\mu_a p_C + \mu_b p_D + \mu_c p_A} + \frac{(\sigma_a p_C + \sigma_b p_D + \sigma_c p_A)^2}{2(\mu_a p_C + \mu_b p_D + \mu_c p_A)^2}$$

$$\sum_i \left(\frac{w_i}{a_i p_C + b_i p_D + c_i p_A} \right) = \frac{\sum w_i}{\mu_a p_C + \mu_b p_D + \mu_c p_A} + \frac{(\sigma_a p_C + \sigma_b p_D + \sigma_c p_A)^2}{2(\mu_a p_C + \mu_b p_D + \mu_c p_A)^2} \sum w_i$$

approximation used in deriving limitation plane

approximation error